

Homework 1

ZhuQinlin 2015K8009929017

September 9, 2017

1 Answer Sheet

Theorem 1 *In the process of graph search, define the set of visited nodes as V , unexplored nodes as U and the fringe as F . For any node $v \in V$, and any node $u \in U$, if there is a path from v to u , the path must contain a node $f \in F$.*

Proof 1 *Prove by contradiction. Assume that there is a path from u to v without any node in F . By the definition of fringe, for any node in V , its adjacent node is either in V , or in F . If there is no node in F in the path from v to u , then we can find at least one pair of nodes s and t in the path, where s is in V and t in U , are adjacent. Thus s has an adjacent node in neither F nor V , which is a contradiction. So the assumption is wrong.*

Theorem 2 *Uniform-cost search gives out the optimal solution, that is, the solution in which the sum of all the actions' cost is minimal.*

Proof 2 *We claim that, for every node, the optimal path for it has been found when it is added to the explored set. This can be proved by contradiction and induction.*

BS: The first node added into the explored set is the initial state, whose optimal solution as a goal state is obviously taking no action, costing zero.

IH: The optimal solutions for all the nodes in the explored set have been found.

IS: If the explored set is consisted of nodes whose optimal solution have been found and the set's fringe has also been constructed, then the lowest-cost node f in the fringe has also found its optimal solution. Because if the real optimal path p only consists of nodes in explored set and f , it is exactly this found solution, or at least as long as it. If not, the path must includes other nodes f' in the fringe, f can't be the lowest-cost node. So by contradiction, the optimal solution from initial state to f has been found.

By induction, the claim above is true. And since finally the goal state is added to the explored set, its optimal solution is found.